Lec 20: Computation Expression

CS220: Programming Principles

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Computation Expression
Recall Asynchronous Computations

async { exprs ... } was an example of computation expressions. Computation expressions provide a convenient syntax for writing computations.
Another Example: `seq`

The `seq` computation expression has a similar form: `seq { ... }`. And it helps build sequence expressions.

```
seq { for i = 0 to 5 do yield (i, i * i) }
seq { while true do yield 1 } // Infinite sequence.
```

Compared to `Seq.unfold`, which one is easier to understand?
What is Similar?

Given the above examples of computation expressions, i.e., seq and async, is there any common thing that you can figure out?
What is Similar?

Given the above examples of computation expressions, i.e., seq and async, is there any common thing that you can figure out?

Those expressions represent specific computations under a specific context.
Expressing Context-Sensitive Computations

We can represent context-sensitive computations using a “wrapped” type\(^1\), called *computation*.

\(^1\)They are wrapped by a type constructor.
Attendance Check

Note:

1. This slide appears at random time during the class.
2. This link is only valid for a few minutes.
3. We don’t accept late responses.
Monads

A monad is a design pattern that allows structuring programs generically while automating away boilerplate code needed by the program logic\(^2\).

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\(^2\)https://en.wikipedia.org/wiki/Monad_(functional_programming)
What Were Hidden in `async` and `seq`?

- When building an `async` computation, all the low-level thread management code was hidden from the expression.
- When building a `seq` computation, the details about constructing cons cells and maintaining its state were hidden.
Monads in this Class

A monad is a triple of a generic type constructor \( M<'T> \) and the two following operations.

\[
\text{val Bind: } M<'T> \times ('T \to M<'U>) \to M<'U>
\]

\[
\text{val Return: } 'T \to M<'T>
\]

✓ Monads are like a box: it is a box with a value of type \( 'T \), and the function \( ('T \to M<'U>) \) takes the value from the box and returns another box of type \( M<'U> \).
Motivating Example

```plaintext
let inc x = x + 1
let dec x = x - 1
// Concise and easy to understand
let id = inc >> dec
```

Function composition is elegant and easy to understand.
Motivating Example (cont’d)

```ocaml
type ResultWithDebugMessage<'a> = {
    Result: 'a
    DbgMsg: string
}
let inc x = { Result = x + 1; DbgMsg = "incremented" }
let dec x = { Result = x - 1; DbgMsg = "decremented" }
let id = inc >> dec // type mismatch.

let id x =
    let rA = inc x
    let rB = dec rA.Result
    { Result = rB.Result; DbgMsg = rA.DbgMsg + "\n" + rB.DbgMsg }
```

Works, but not elegant.
let bind f r =
  let r' = f r.Result
  { r' with DbgMsg = r.DbgMsg + "\n" + r'.DbgMsg }

let id = inc >> bind dec

The id function now combines inc and dec in an elegant manner.
Return the Boxed Type

We want to make `ResultWithDebugMessage` value from an integer: `wrap` function simply wraps a value without any debugging message.

```
let wrap r = { Result = r; DbgMsg = "" }
let id =
  inc
  >> bind dec
  >> bind (fun x -> x + 1 |> wrap)
  >> bind dec
```
Signatures of `bind` and `wrap`

val bind:  
  ('a -> ResultWithDebugMessage<'b>)  
  -> ResultWithDebugMessage<'a>  
  -> ResultWithDebugMessage<'b>

val wrap:  
  'a -> ResultWithDebugMessage<'a>

val Bind:  
  M<'T> * ('T -> M<'U>) -> M<'U>

(Same as `bind` after swapping the argument order)

val Return:  
  'T -> M<'T>

ResultWithDebugMessage<'a> was an example of monad!
Bottom Line

We can logically bind functions while hiding some details behind the scene with monads. Typically we define a bind operator (\(\gg=\)) which is an infix-operator for the Bind function discussed above.

```ocaml
let (\(\gg=\)) m f = bind f m
let id x = inc x \(\gg=\) dec \(\gg=\) inc \(\gg=\) dec
```
Enhancing The Expressiveness
Let-Bindings Revisited

We can always convert let-bindings to a function with nested function calls: function calls another function, and the function calls another function, and so on.

```
let x = 1
let y = 2 + x
let z = x * y
z
```

```
1 |> fun x ->
  2 + x |> fun y ->
    x * y |> fun z ->
      z
```
Creating a Bind Function

Let us now create a `bind` function that takes in a value and a function, and apply the value to the function (as in the pipe operator).

```ml
let bind x f = f x
let ret x = x
bind 1 (fun x ->
    bind (2 + x) (fun y ->
        bind (x * y) (fun z ->
            ret z)))
```

This is so-called "continuation passing style" (CPS).
Continuation Passing Style?

In CPS, functions always end with a function that we call \textit{continuation}, which describes what to do next.

\begin{verbatim}
// Normal
let add a b = a + b
let x = add (add 1 2) 3
// CPS
let add a b cont = cont (a + b)
let x = add 1 2 (fun r -> add r 3 (fun r -> r))
\end{verbatim}

\textbf{★ Food for Thought.} CPS forces us to write tail-recursive functions, but it doesn’t mean that it helps reduce memory consumption.
CPS Bindings vs. Let-Bindings

Syntactically different but semantically the same.

```
bind 1 (fun x ->
    bind (2 + x) (fun y ->
        bind (x * y) (fun z ->
            ret z)))

let x = 1
let y = 2 + x
let z = x * y
z
```
Hiding Complex Logic

If we can transform the chain of `bind-ret` function calls into expressions that look like `let-bindings`, and if our language supports such a transformation, then we can hide some complex logic under a beautiful language.

```ml
let bind x f =
  printfn "you can do some complex things here."
  f x

bind 1 (fun x ->
  bind (2 + x) (fun y ->
    bind (x * y) (fun z ->
      ret z)))
```
let vs. let!

Compare bind and ResultWithDebugMessage.bind

```ocaml
let bind x f = f x // 'a -> ('a -> 'b) -> 'b

module ResultWithDebugMessage = 
  // ResultWithDebugMessage<'a> ->
  // ('a -> ResultWithDebugMessage<'b>) ->
  // ResultWithDebugMessage<'b>
  let bind r f =
    let r' = f r.Result
    { r' with DbgMsg = r.DbgMsg + "\n" + r'.DbgMsg }
```
Example: Safe Division

Safe division function.

```plaintext
let safeDiv a b =
  if b = 0 then None
else Some (a / b)
```
Too Many Nested Checks

```plaintext
// unsafe div
let x = (((a / b) / c) / d) / e

// safe div
let x' =
    match safeDiv a b with
    | None -> None
    | Some r ->
        match safeDiv r c with
        | None -> None
        | Some r ->
            match safeDiv r d with
            | None -> None
            | Some r -> safeDiv r e
```
Observation

The nested match expressions follow a CPS. We can write our own “bind” function to connect them!

```ocaml
let bind (x, f) = 
  match x with 
  | None -> None 
  | Some m -> f m

let ret x = Some x
```

```ocaml
bind (safeDiv a b, fun r -> 
  bind (safeDiv r c, fun r -> 
    bind (safeDiv r d, fun r -> 
      bind (safeDiv r e, fun r -> 
        ret r))))
```
Built-in Binder: Option.bind

Works the same, but it takes a continuation first.

```ocaml
val Option.bind : ('a -> 'b option) -> 'a option -> 'b option

safeDiv a b |> Option.bind (fun r ->
    safeDiv r c |> Option.bind (fun r ->
        safeDiv r d |> Option.bind (fun r ->
            safeDiv r e |> Option.bind (fun r ->
                ret r))))
```
Computation Expression Builder

A computation expression builder is a class that contains several member functions such as Bind and Return.

val __.Bind:  M<'T> * ('T -> M<'U>) -> M<'U>
val __.Return:  'T -> M<'T>

We can define our own builder to create our computation expressions.
Maybe Computation Expression

type MaybeBuilder () =
  member __.Bind (m, f) = Option.bind f m
  member __.Return (m) = Some m

let maybe = MaybeBuilder ()
maybe {  
    let! r = safeDiv a b  
    let! r = safeDiv r c  
    let! r = safeDiv r d  
    let! r = safeDiv r e  
    return r  
}

- The **Bind** member corresponds to the `let!` expression.  
- The **Return** member corresponds to the `return` expression.
match Email.create emailInput with
| Some email ->
  match Name.create nameInput with
  | Some name ->
    match ID.create idInput with
    | Some id ->
      DB.insert db email name id // insert into the DB
    | None -> db // return the DB as it is
  | None -> db
| None -> db
DB Example (cont’d)

```ml
maybe {
  let! email = Email.create emailInput
  let! name = Name.create nameInput
  let! id = ID.create idInput
  return DB.insert email name id
}
```

Elegant and easy to read!
In-Class Activity #20
Preparation

We are going to use the same git repository as before. Just in case you don’t have it, clone the repository using the following command.

1. Clone the repository to your machine.

2. Move in to the directory CS220-Main/Activities
   - cd CS220-Main
   - cd Activities
Problem: Define Your Own

Define your own computation expression for list. The computation expression should be able to handle the following code.

```
mylist { for i in [1 .. 10] do yield i * i }
mylist { for i in [1 .. 10] do yield i * i }
```

Hint: you need to define Yield and For members.
for identifier = enumerable-expr do cexpr

becomes

builder.For (enumerable-expr, fun identifier -> cexpr)

for identifier = expr1 to expr2 do cexpr

becomes

builder.For (seq [expr1 .. expr2], fun identifier -> cexpr)
Conclusion
Further Readings

- https://fsharpforfunandprofit.com/series/computation-expressions.html
Question?